Computer Organization and Structure

Homework #5 Due: 2007/12/25

1. Consider the single-cycle datapath in Figure 1. A friend is proposing to modify this single-cycle datapath by eliminating the control signal MemtoReg. The multiplexor that has MemtoReg as an input will instead use either the ALUSrc or the MemRead control signal. Will your friend's modification work? Can one of the two signals (MemRead and ALUSrc) substitute for the other? Explain. Furthermore, determine whether any of the control signals in the single-cycle implementation can be eliminated and replaced by another existing control signal, or its inverse. Note that such redundancy is there because we have a very small set of instructions at this point, and it will disappear (or be harder to find) when we implement a larger number of instructions.



Figure 1: The simple datapath with the control unit.

2. We wish to add the instructions jr (jump register), sll (shift left logical), lui (load upper immediate), and a variant of the lw (load word) instruction to the single-cycle datapath. The variant of the lw instruction increments the index register after loading word from memory. This instruction (l_inc) corresponds to the following two instructions:

| lw | \$rs, | L(\$rt) | | |
|------|-------|---------|--|--|
| addi | \$rt, | \$rt, 1 | | |

Add any necessary datapaths and control signals to Figure 1 and show the necessary additions to Table 1. You can photocopy Figure 1 and Table 1 to make it faster to show the additions.

| | | | Memto | Reg | Mem | Mem | | | |
|-------------|--------|--------|-------|-------|------|-------|--------|--------|--------|
| Instruction | RegDst | ALUSrc | Reg | Write | Read | Write | Branch | ALUOp1 | ALUOp0 |
| R-format | 1 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 |
| lw | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 |
| SW | Х | 1 | X | 0 | 0 | 1 | 0 | 0 | 0 |
| beq | Х | 0 | X | 0 | 0 | 0 | 1 | 0 | 1 |

Table 1: The setting of the control lines is completely determined by the opcode fields of the instruction.

3. Two important parameters control the performance of a processor: cycle time and cycles per instruction. There is an enduring trade-off between these two parameters in the design process of microprocessors. While some designers prefer to increase the processor frequency at the expense of large CPI, other designers follow a different school of thought in which reducing the CPI comes at the expense of lower processor frequency. Consider the following machines, and compare their performance using the following instruction mix: 25% loads, 13% stores, 47% ALU instructions, and 15% branches/jumps.

M1: The multicycle datapath is designed as shown in Figure 3 with a 1 GHz clock.

M2: A machine like M1 except that register updates are done in the same clock cycle as a memory read of ALU operation. Thus in Figure 4, states 6 and 7 and states 3 and 4 are combined. This machine has an 3.2 GHz clock, since the register update increases the length of the critical path.

M3: A machine like M2 except that effective address calculations are done in the same clock cycle as a memory access. Thus states 2, 3, and 4 can be combined, as can 2 and 5, as well as 6 and 7. This machine has a 2.8 GHz clock because of the long cycle created by combining address calculation and memory access.

Find out which of the machines is fastest. Are there instruction mixes that would make another machine faster, and if so, what are they?



Figure 2: The multicycle datapath with the control lines.



Figure 3: The complete datapath for the multicycle implementation together with the necessary control lines.



Figure 4: The complete finite state machine control for the datapath shown in Figure 2.