Computer Organization and Structure

Homework #5 Due: 2007/1/9

1. If the time for an ALU operation can be shortened by 25% (compared to the following table);

| Instruction class | Instruction | Register | ALU | Data | Register | Total |
|---------------------|-------------|----------|-----------|--------|----------|--------|
| | fetch | read | operation | access | write | time |
| Load word (lw) | 200 ps | 100 ps | 200 ps | 200 ps | 100 ps | 800 ps |
| Store word (sw) | 200 ps | 100 ps | 200 ps | 200 ps | | 700 ps |
| R-format (add, sub, | 200 mg | 100 ps | 200 ps | | 100 ps | 600 ps |
| and, or, slt) | 200 ps | | | | | |
| Branch (beq) | 200 ps | 100 ps | 200 ps | | | 500 ps |

- a. Will it affect the speedup obtained from pipelining? If yes, by how much? Otherwise, why?
- b. What if the ALU operation now takes 25% more time?
- 2. Identify all of the data dependencies in the following code. Which dependencies are data hazards that will be resolved via forwarding? Which dependencies are data hazards are will cause a stall?

| add | \$3, | \$4, | \$2 |
|-----|------|------|------|
| sub | \$5, | \$3, | \$1 |
| lw | \$6, | 200(| \$3) |
| add | \$7, | \$3, | \$6 |

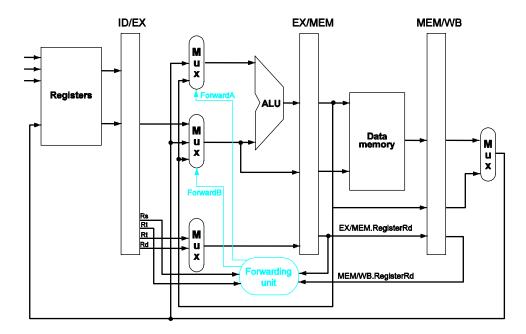
3. The following piece of code is executed using the pipeline shown in the following figure:

| \$5, | 40(\$2) |
|------|----------------------|
| \$6, | \$3, \$2 |
| \$7, | \$2, \$1 |
| \$8, | \$4, \$3 |
| \$9, | \$2, \$1 |
| | \$6, \$7, \$8, |

At cycle 5, right before the instructions are executed, the processor state is as follows:

- a. The PC has the value 100_{ten}, the address of the sub_instruction.
- b. Every register has the initial value 10_{ten} plus the register number (e.g., register \$8 has the initial value 18_{ten}).
- c. Every memory word accessed as data has the initial value 1000_{ten} plus the byte address of the word (e.g., Memory[8] has the initial value 1008_{ten}).

Determine the value of every field in the four pipeline registers in cycle 5.



4. The performance can be *maximized* on the pipelined datapath by using forwarding and stalls on a use following a load. Rewrite the following code to *minimize* performance on this datapath – that is, reorder the instructions so that this sequence takes the *most* clock cycles to execute while stall obtaining the same result.

| lw | \$2, | 100(| \$6) |
|-----|------|------|------|
| lw | \$3, | 200(| \$7) |
| add | \$4, | \$2, | \$2 |
| add | \$6, | \$3, | \$5 |
| sub | \$8, | \$4, | \$6 |
| lw | \$7, | 300(| \$8) |
| beq | \$7, | \$8, | Loop |
| | | | |