

# *Surface Simplification Using Quadric Error Metrics (QEM)*

Michael Garland

Paul Heckbert

SIGGRAPH 97

***Demo***

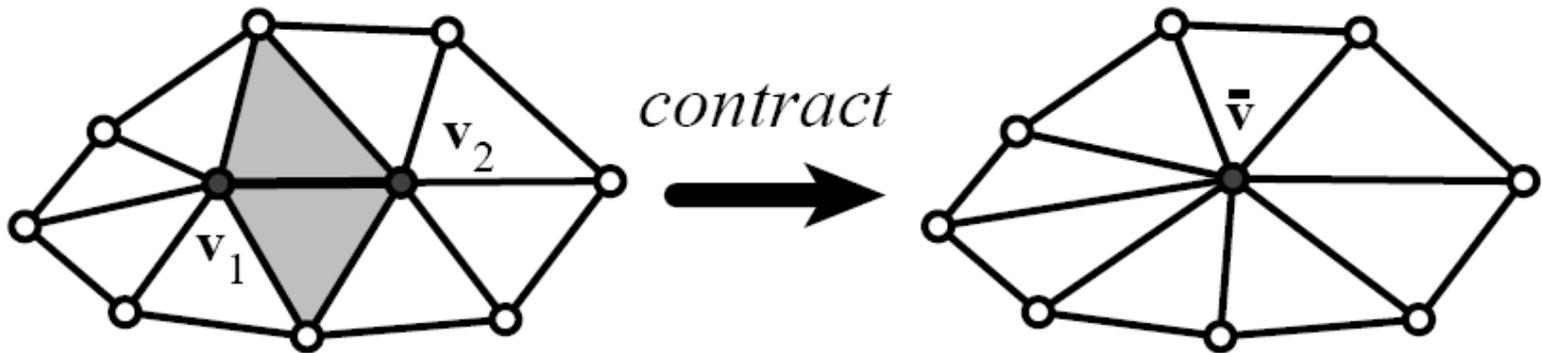
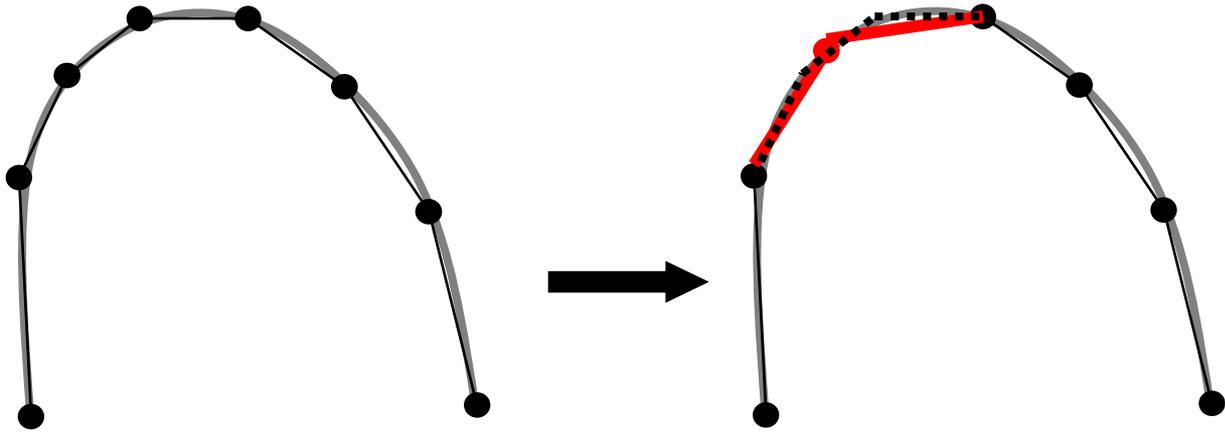
# Objectives

- Fast and good quality
  - Efficiency and quality trade-off
  
- Convenient characterization of error/shape
  - Compact and efficient to compute

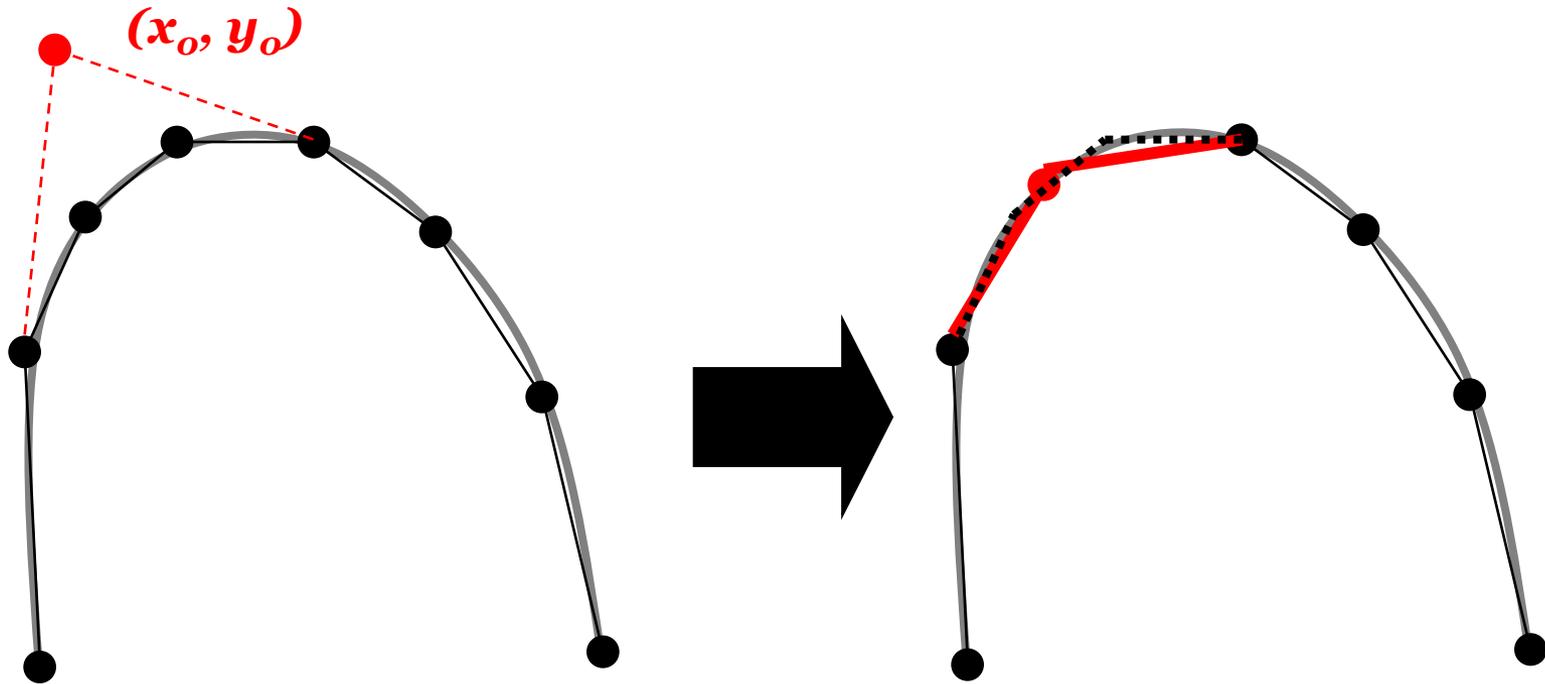
# Things to consider

- 1. Defining the basic operation**
- 2. Defining the error metrics**
3. Primitives removal strategy
4. Remaining updates

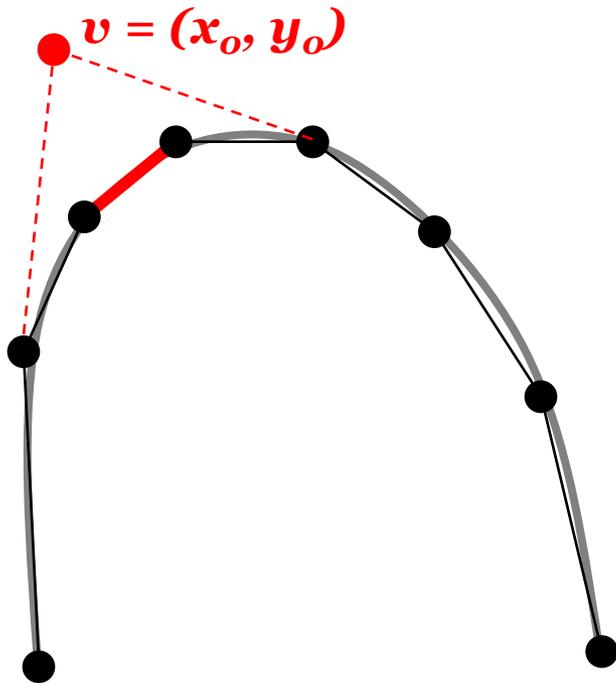
# Basic operation



# The error metrics – 2D



# The error metrics – 2D



*Matrix form*

$$\begin{bmatrix} a_1 & b_1 & c_1 \end{bmatrix} * \begin{bmatrix} x_o \\ y_o \\ 1 \end{bmatrix}$$

$$L_1 = \begin{bmatrix} a_1 \\ b_1 \\ c_1 \end{bmatrix} \quad v = \begin{bmatrix} x_o \\ y_o \\ 1 \end{bmatrix}$$

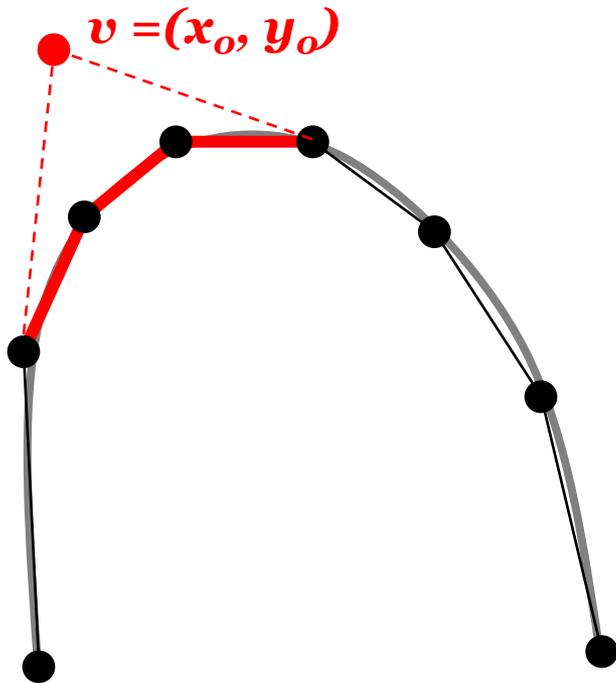
$$d_1^2 = \begin{bmatrix} x_o & y_o & 1 \end{bmatrix} * \begin{bmatrix} a_1 \\ b_1 \\ c_1 \end{bmatrix} * \begin{bmatrix} a_1 & b_1 & c_1 \\ x_o \\ y_o \\ 1 \end{bmatrix}$$

$$d_1^2 = v^T * (L_1 * L_1^T) * v = v^T * (Q_1) * v$$

*Line segment  $L_1$ :  $a_1x + b_1y + c_1 = 0$*

*$(x_o, y_o)$  to  $L_1$  is  $d_1 = |a_1 * x_o + b_1 * y_o + c_1|$*

# The error metrics – 2D



**Total** squared distance

$$E_v = \sum_i v^T * (L_i * L_i^T) * v = v^T * (\sum_i Q_i) * v$$

To better fit the curve  
minimize squared distance

$$dE_v/dv = (\sum_i Q_i) * v = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}$$

So best  $v$  is at  $(\sum_i Q_i)^{-1} * \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}$

# For meshes

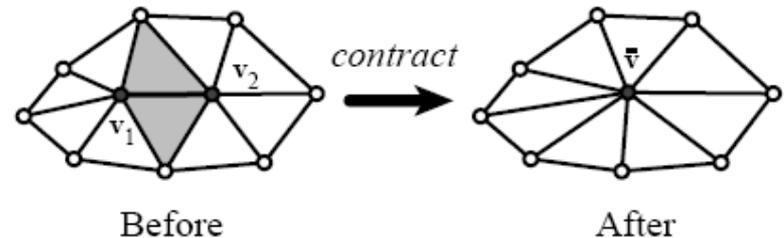
- Use plane equation instead of line equation

- $ax + by + cz + d = 0$

- $a^2 + b^2 + c^2 = 1$

- $v = [v_x \ v_y \ v_z \ 1]^T$

- $p = [a \ b \ c \ d]^T$



$$\Delta(\mathbf{v}) = \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} (\mathbf{v}^T \mathbf{p})(\mathbf{p}^T \mathbf{v})$$

$$= \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} \mathbf{v}^T (\mathbf{p}\mathbf{p}^T) \mathbf{v}$$

$$= \mathbf{v}^T \left( \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} \mathbf{K}_p \right) \mathbf{v}$$

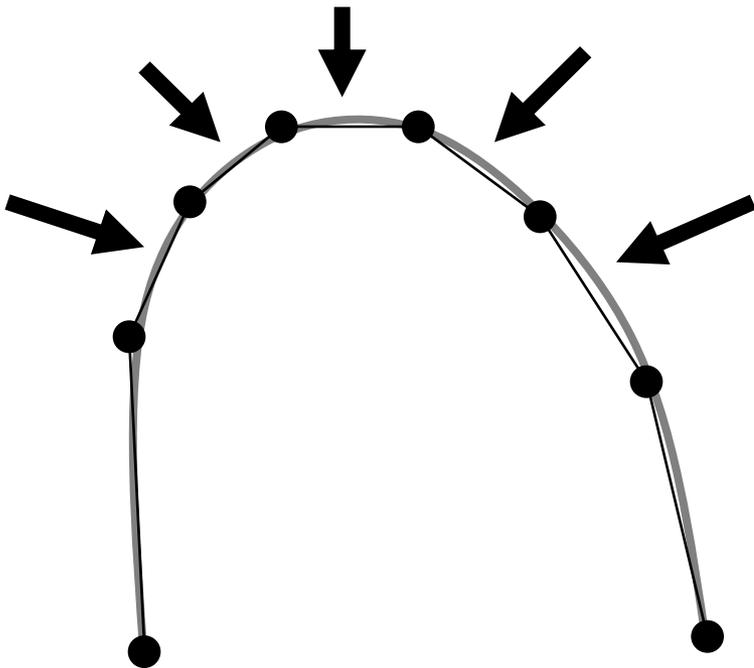
$$\mathbf{K}_p = \mathbf{p}\mathbf{p}^T = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$

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# Primitives removal strategy

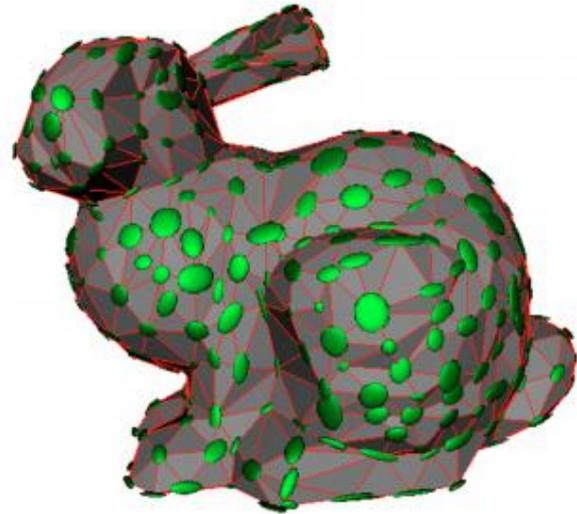
*But which one?*



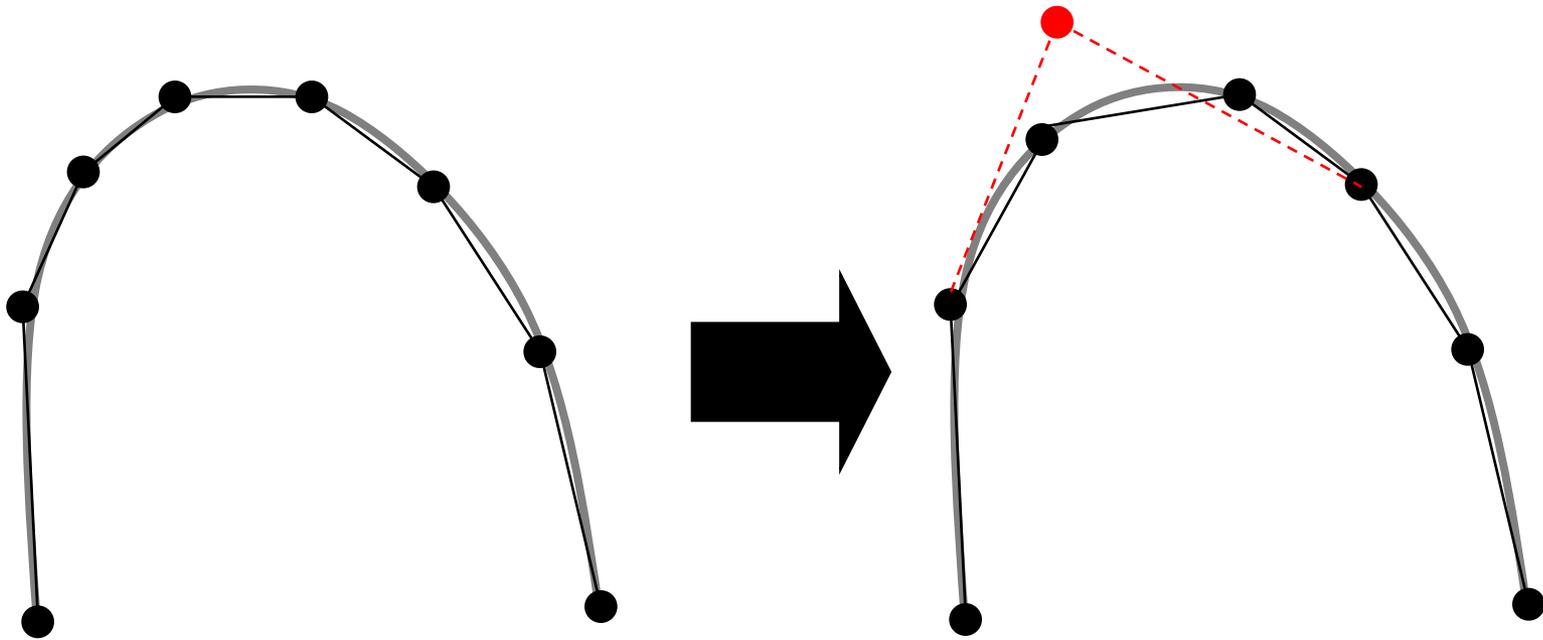
Total squared distance

$$E_v = \sum_i v^T * (L_i * L_i^T) * v = v^T * (\sum_i Q_i) * v$$

Use priority queue!



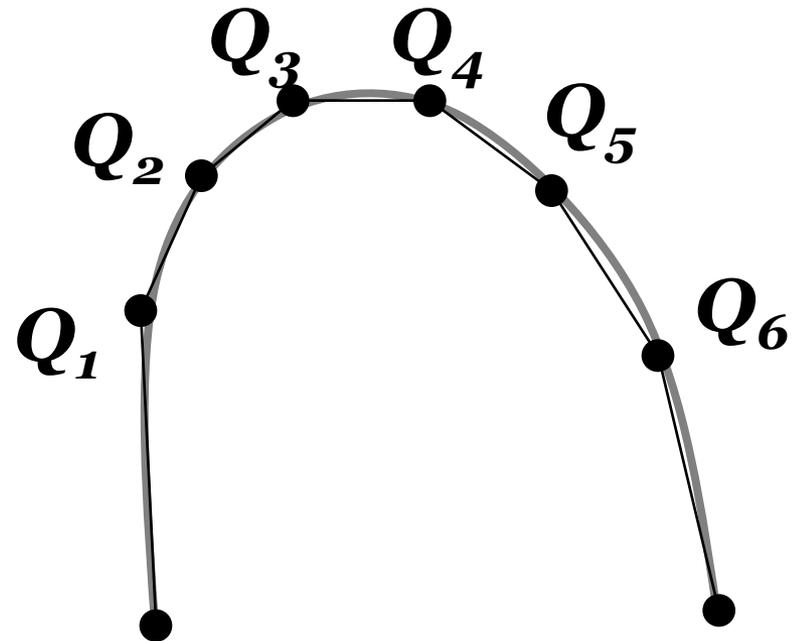
# Efficiency consideration



- *Accumulated error w/ respect to **ALL** underlying segments!*
- *How many do we have in queue?*

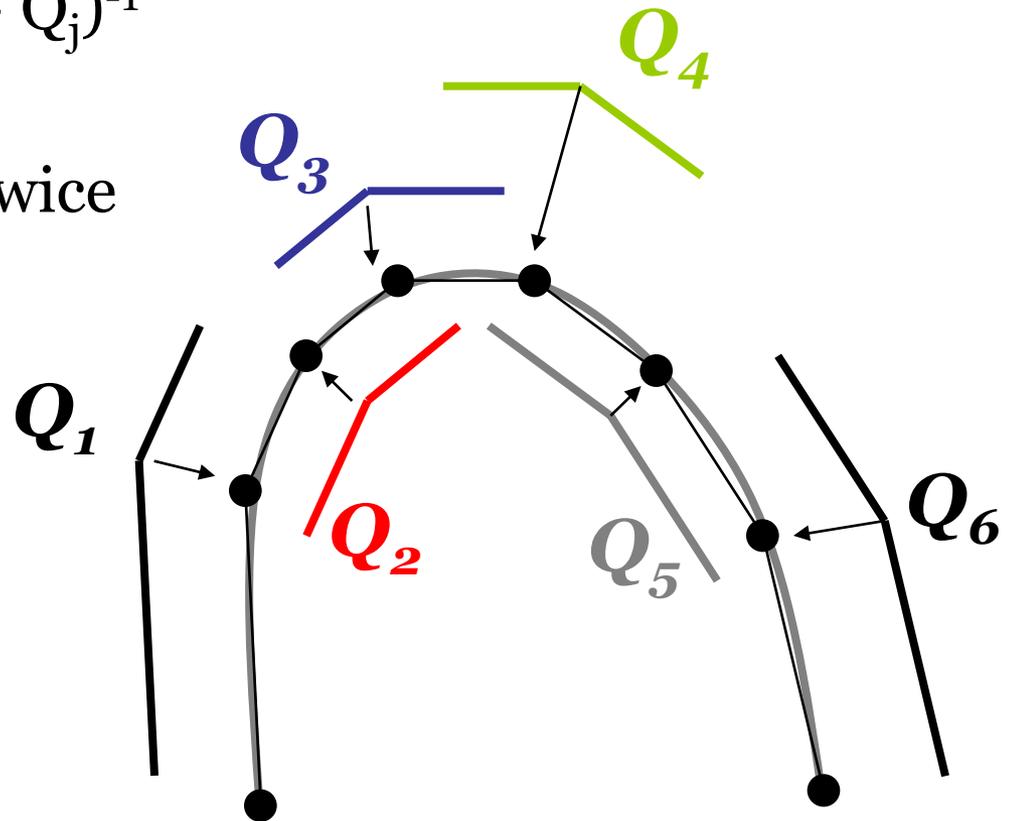
# Cheats!

- Keep indices to all underlying segments!
  - Storage and indexing speed.
  - Sometimes you only compare errors
- So vertices remember their own incident segments as  $Q$ 
  - $V_{\text{new}} = \text{minimize } (Q_i + Q_j)^{-1}$
  - Just sum two
  - The middle counted twice
  - For meshes, 3 times



# Cheats!

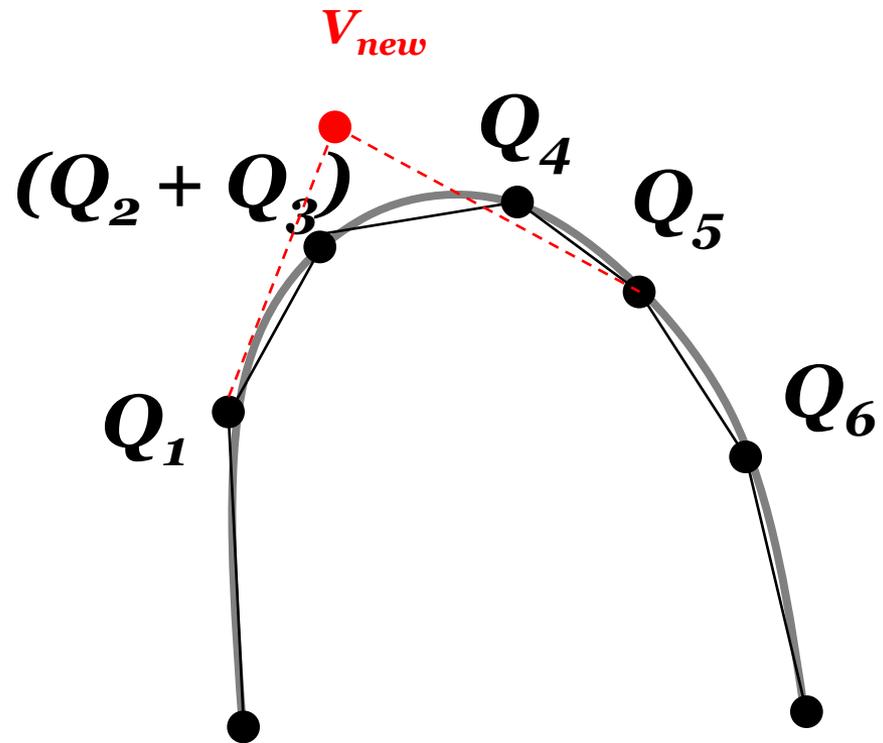
- So vertices remember their own incident segments as  $Q$ 
  - $V_{new} = \text{minimize } (Q_i + Q_j)^{-1}$
  - Just sum two
  - The middle counted twice



# Remaining Updates

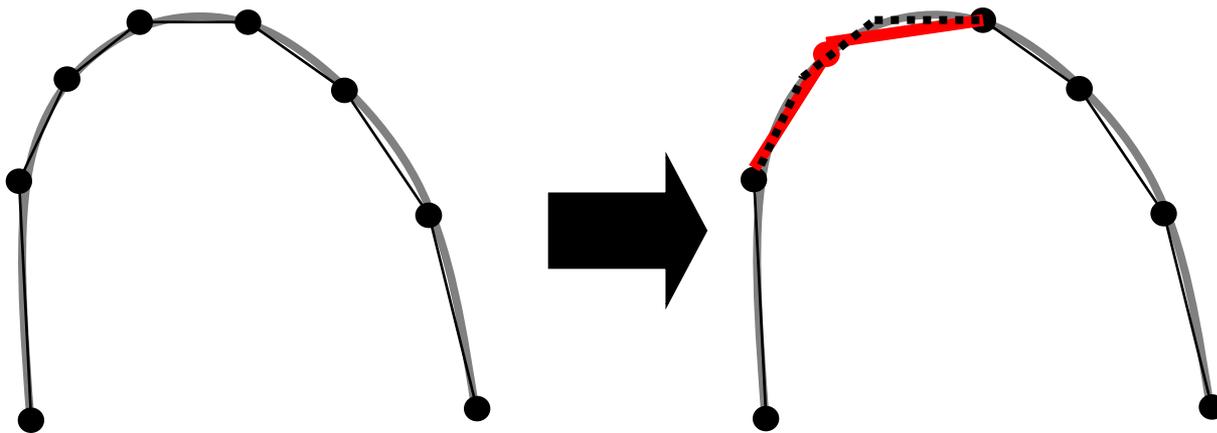
• **Error** ( $V_{new}$ ) =  
$$V_{new}^T (Q_2 + Q_3 + Q_4) V_{new}$$

• **Minimized error to find**  $V_{new}$



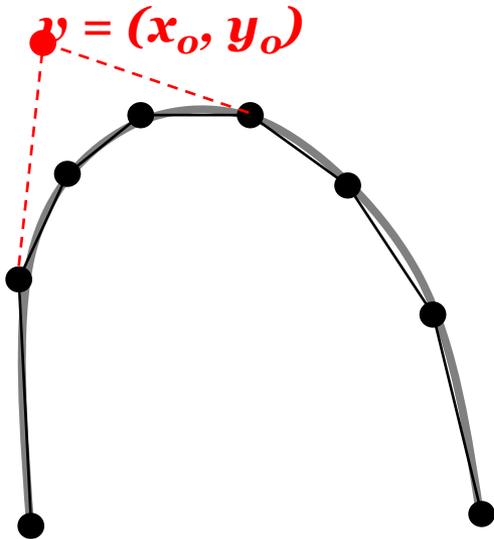
# Quick review

- 1. Defining the basic operation**
2. Defining the error metrics
3. Primitives removal strategy
4. Remaining updates



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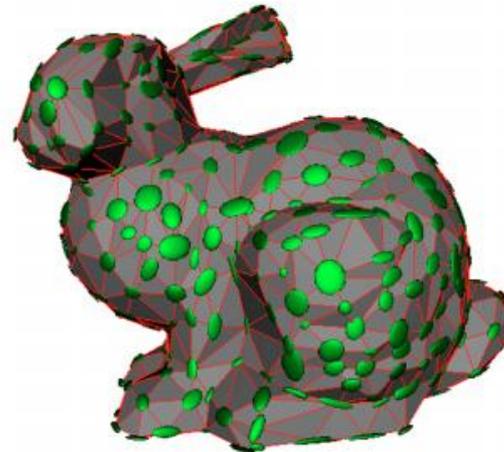
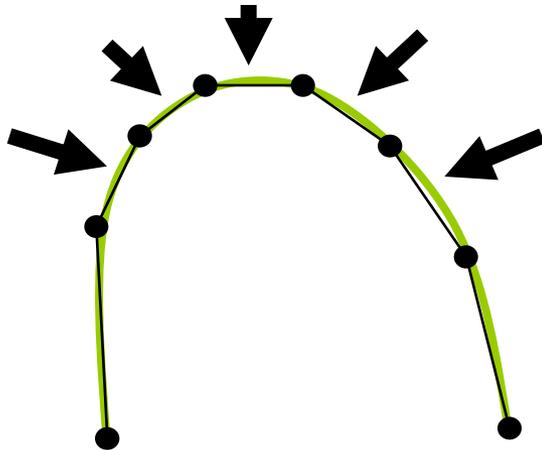
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$$\begin{aligned}\Delta(\mathbf{v}) &= \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} (\mathbf{v}^T \mathbf{p})(\mathbf{p}^T \mathbf{v}) \\ &= \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} \mathbf{v}^T (\mathbf{p} \mathbf{p}^T) \mathbf{v} \\ &= \mathbf{v}^T \left( \sum_{\mathbf{p} \in \text{planes}(\mathbf{v})} \mathbf{K}_{\mathbf{p}} \right) \mathbf{v}\end{aligned}$$

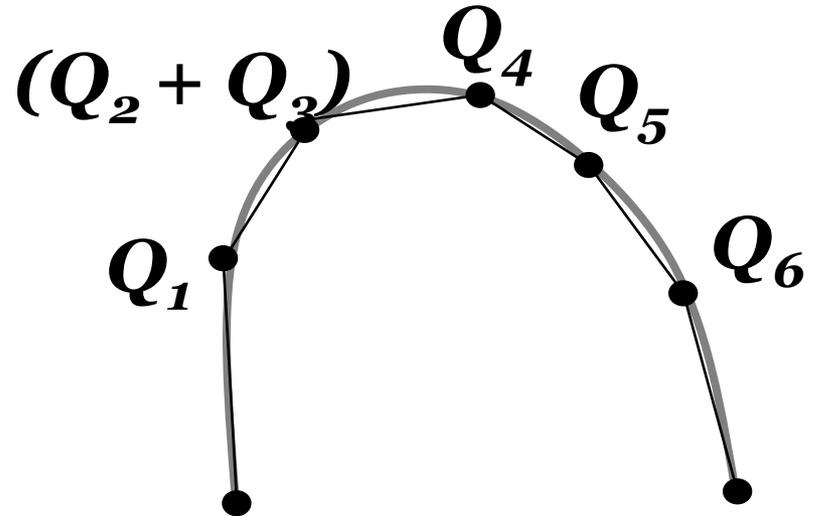
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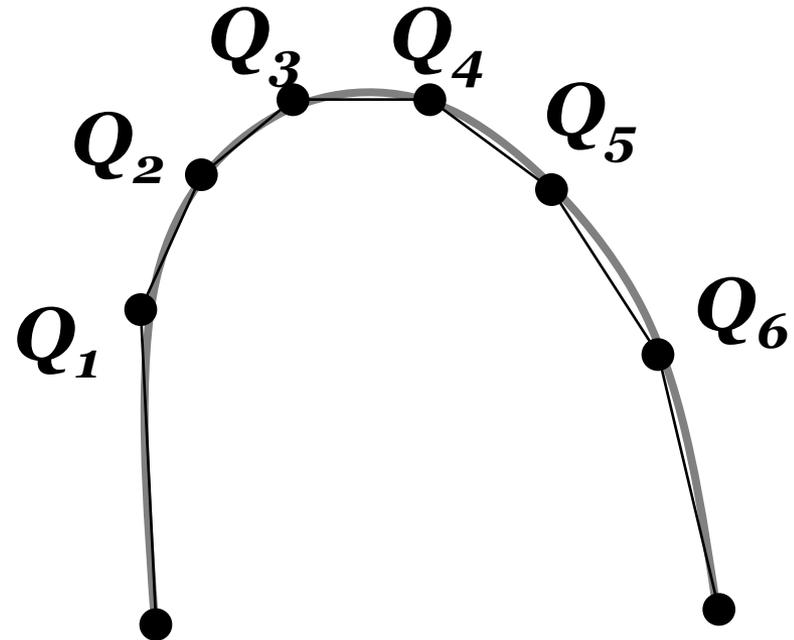
# Quick review

1. Defining the basic operation
2. Defining the error metrics
3. Primitives removal strategy
4. **Remaining updates**



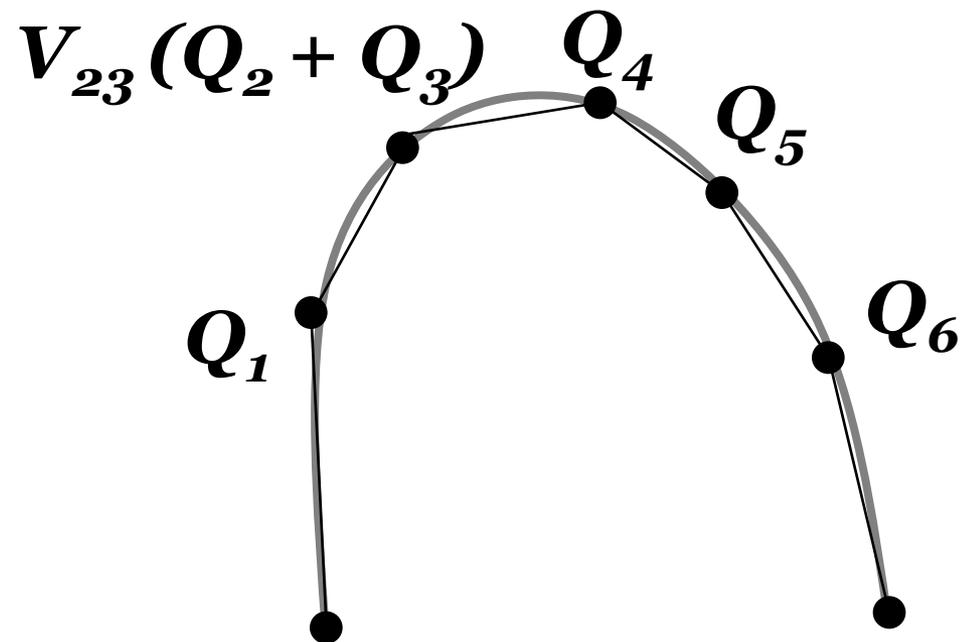
# Illustration of operations

1. Initial  $Q$  for every vertices
2. Compute  $v_{\text{new}}$  and  $\text{err}(v_{\text{new}})$  for every edges
  1.  $V_{1-2}, V_{2-3}, V_{3-4}, V_{4-5} \dots$
  2.  $(Q_1+Q_2)^{-1}, (Q_2+Q_3)^{-1}, (Q_3+Q_4)^{-1}, (Q_4+Q_5)^{-1} \dots$
  3. Put into min heap



# Illustration of operations

3. Select  $\min \text{err}(V_{ij})$  to contract
4. Update **necessary items** in the min heap
  1.  $V_{1-23}, V_{23-4}, V_{3-4}, V_{4-5}$
  2. Their  $Q$
5. Repeat step-3



# Iterations from paper

1. Compute the  $Q$  matrices for all the initial vertices
2. Compute the optimal contraction target  $\mathbf{v}_{\text{new}}$  for each valid pair  $\mathbf{v}_1 \mathbf{v}_2$ . The error  $\mathbf{v}_{\text{new}}^T(\mathbf{Q}_1 + \mathbf{Q}_2)\mathbf{v}_{\text{new}}$  of this target vertex becomes the *cost* of contracting that pair.
3. Place all the pairs in a heap keyed on cost with the minimum cost pair at the top.
4. Iteratively remove the pair  $\mathbf{v}_1 \mathbf{v}_2$  of least cost from the heap, contract this pair, and update the costs of all valid pairs involving  $\mathbf{v}_{\text{New}}$

# Objectives

- Fast and good quality
  - Efficiency and quality trade-off
  - No indexing
  - Only require necessary updates
- Convenient characterization of error/shape
  - Compact and efficient to compute
  - 10 coefficient, 4x4 matrix inversion/addition

# Something not addressed...

- Focused on 2-manifold
  - no joining
- Boundary condition
  - Eaten away
- **Triangle inversion**
  - Why have the case?

